

# Comp 115: Databases

## The Relational Model

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<http://www.cs.tufts.edu/comp/115/>

# The Relational Model

Intro & SQL overview

Keys & Integrity Constraints

ER to Relational

ISA to Relational

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# Why the Relational Model?

most widely used model

*IBM, Microsoft, Oracle, etc.*

”Legacy systems” in older models

e.g., IBM’s IMS

object-relational model incorporates oo concepts

*IBM DB2, Oracle 11i*

more recently: key-value store

# Relational

tables with rows and columns

well-defined schema

data model fits data rather than functionality

deduplication

# Key/Value

collections of documents

schema-less (each document can have different schema)

data stored in an application-friendly way

possible duplication

# Relational Database: Definitions

*relational database*: a collection (set) of *relations*

*each relation*: made up of 2 parts

*schema*: name of relation, name & type of each column

Students(*sid*: string, *name*: string, *login*: string, *age*: integer, *gpa*: real)

*instance* : a *table*, with rows and columns.

#rows = *cardinality*

#fields = *degree / arity*

can think of a relation as a *set* of rows or *tuples*

- (1) all rows are distinct
- (2) no order among rows

# Instance of Students Relation

sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@cs	18	3.2
53650	Smith	smith@math	19	3.8

cardinality = 3, arity = 5, all rows distinct

do all values in each column of a relation instance have to be distinct?



# SQL - A language for Relational DBs

SQL\* (a.k.a. “Sequel”), standard language

## Data Definition Language (DDL)

create, modify, delete relations

specify constraints

administer users, security, etc.

## Data Manipulation Language (DML)

specify *queries* to find tuples that satisfy criteria

add, modify, remove tuples

\* Structured Query Language



# SQL Overview

```
CREATE TABLE <name> ( <field> <domain>, ... )
```

```
INSERT INTO <name> (<field names>)  
VALUES (<field values>)
```

```
DELETE FROM <name>  
WHERE <condition>
```

```
UPDATE <name>  
SET <field name> = <value>  
WHERE <condition>
```

```
SELECT <fields>  
FROM <name>  
WHERE <condition>
```

# Creating Relations in SQL

type (**domain**) of each field is specified

also enforced whenever tuples are added or modified

```
CREATE TABLE Students  
  (sid CHAR(20),  
   name CHAR(20),  
   login CHAR(10),  
   age INTEGER,  
   gpa FLOAT)
```

# Table Creation (continued)

Enrolled: holds information about courses  
students take

```
CREATE TABLE Enrolled  
  (sid CHAR(20),  
   cid CHAR(20),  
   grade CHAR(2))
```

# Adding and Deleting Tuples

Can insert a single tuple using:

```
INSERT INTO Students (sid, name, login, age, gpa)
VALUES ('53688', 'Smith', 'smith@cs', 18, 3.2)
```

Can delete all tuples satisfying some condition  
(e.g., name = Smith):

```
DELETE
FROM Students S
WHERE S.name = 'Smith'
```

Powerful variants of these commands are available;  
more later!

# The Relational Model

Intro & SQL overview

**Keys & Integrity Constraints**

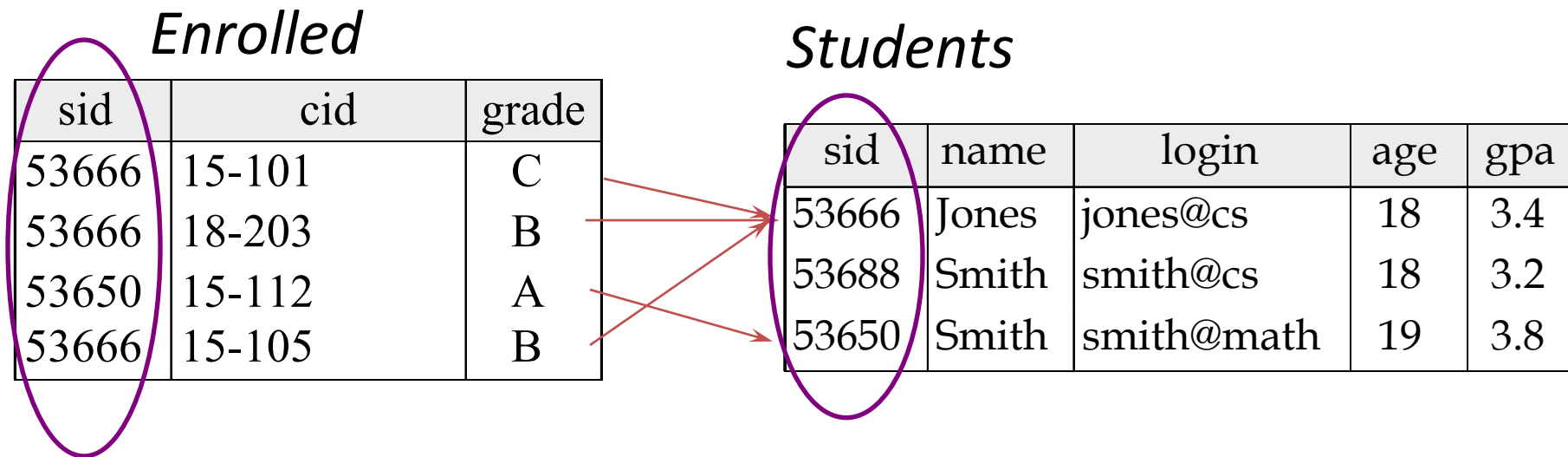
ER to Relational

ISA to Relational

# Keys

keys: associate tuples in different relations

keys are one form of integrity constraint (IC)



**FOREIGN Key**

**PRIMARY Key**

# Primary Keys

sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@cs	18	3.2
53650	Smith	smith@math	19	3.8

A set of fields is a superkey if:

No two distinct tuples can have same values in all key fields

Is <sid> a superkey?

What about <sid,name>?

What about <sid,name,age>?

What about <age,name>?



# Primary Keys

sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@cs	18	3.2
53650	Smith	smith@math	19	3.8

A set of fields is a superkey if:

No two distinct tuples can have same values in all key fields

A set of fields is a key for a relation if :

It is a superkey

No subset of the fields is a superkey



Is  $\langle \text{sid} \rangle$  a key?  ~~$\langle \text{sid}, \text{name} \rangle$~~ ?  ~~$\langle \text{sid}, \text{name}, \text{age} \rangle$~~ ?  ~~$\langle \text{age}, \text{name} \rangle$~~ ?



# Primary Keys

sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@cs	18	3.2
53650	Smith	smith@math	19	3.8

A set of fields is a superkey if:

No two distinct tuples can have same values in all key fields

A set of fields is a key for a relation if :

It is a superkey

No subset of the fields is a superkey



what if >1 key for a relation?

chose one as the **primary key** / rest called **candidate** keys

# Primary and Candidate Keys in SQL

possibly many candidate keys (specified using **UNIQUE**),  
one of which is chosen as the *primary key*

keys must be defined carefully!

"for a given student and course, there is a single grade"

```
CREATE TABLE Enrolled  
(sid CHAR(20)  
  cid CHAR(20),  
  grade CHAR(2),  
  PRIMARY KEY (sid,cid))
```

vs.

```
CREATE TABLE Enrolled  
(sid CHAR(20)  
  cid CHAR(20),  
  grade CHAR(2),  
  PRIMARY KEY (sid),  
  UNIQUE (cid, grade))
```



"students can take only one course, and no two students  
in a course receive the same grade"

# Foreign Keys, Referential Integrity

foreign key: set of fields in one relation that is used to “refer” to a tuple in another  
correspond to the primary key of the other relation  
a “logical pointer”

If all foreign key constraints are enforced, referential integrity is achieved (i.e., no dangling references)

# Foreign Keys in SQL

Example: Only students listed in the Students relation should be allowed to enroll for courses.

*sid* is a foreign key referring to **Students**

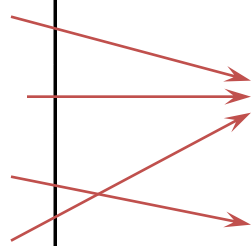
```
CREATE TABLE Enrolled
(sid CHAR(20),cid CHAR(20),grade CHAR(2),
PRIMARY KEY (sid,cid),
FOREIGN KEY (sid) REFERENCES Students )
```

*Enrolled*

sid	cid	grade
53666	15-101	C
53666	18-203	B
53650	15-112	A
53666	15-105	B

*Students*

sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@cs	18	3.2
53650	Smith	smith@math	19	3.8



# Enforcing Referential Integrity

Students and Enrolled; *sid* in Enrolled is a FK references Students

What to do if a tuple with a non-existent *sid* is inserted in Enrolled?

What should be done if a Students tuple is deleted?

Also delete all Enrolled tuples that refer to it?

Disallow deletion of a Students tuple that is referred to?

Set *sid* in Enrolled tuples that refer to it to a *default sid*?

(In SQL we can set *sid* to be equal to *null*, denoting “unknown” or “inapplicable”)

Similar issues arise if primary key of Students tuple is updated



# Integrity Constraints (ICs)

**IC:** must be true for *any* instance of the database  
(e.g., domain constraints)

ICs are specified when schema is defined

ICs are checked when relations are modified

a *legal* instance of a relation satisfies all specified ICs

DBMS should not allow illegal instances

if the DBMS checks ICs, stored data is more faithful to  
real-world meaning  
avoids data entry errors, too!

# Where do ICs Come From?

ICs are based upon the *real-world semantics*

we can check a database instance to see if an IC is violated, but we cannot infer that an IC hold

An IC is a statement about *all possible* instances!

From example, we know *name* is not a key, but the assertion that *sid* is a key is given to us

key and foreign key ICs are the most common  
(more general ICs supported too)

# The Relational Model

Intro & SQL overview

Keys & Integrity Constraints

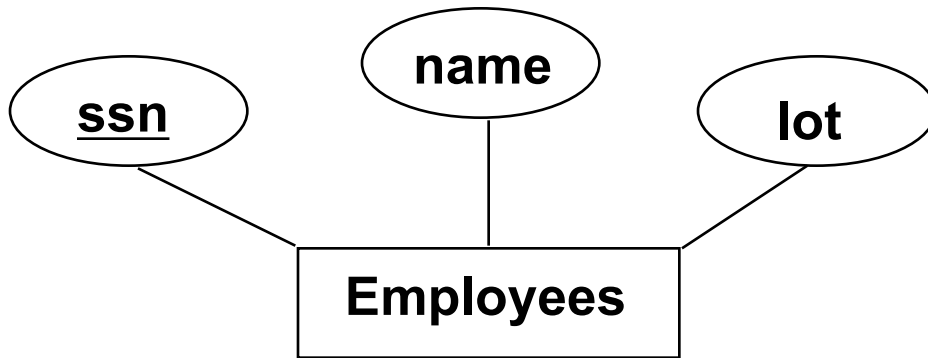
**ER to Relational**

ISA to Relational



# Logical DB Design: ER to Relational

Entity sets to tables

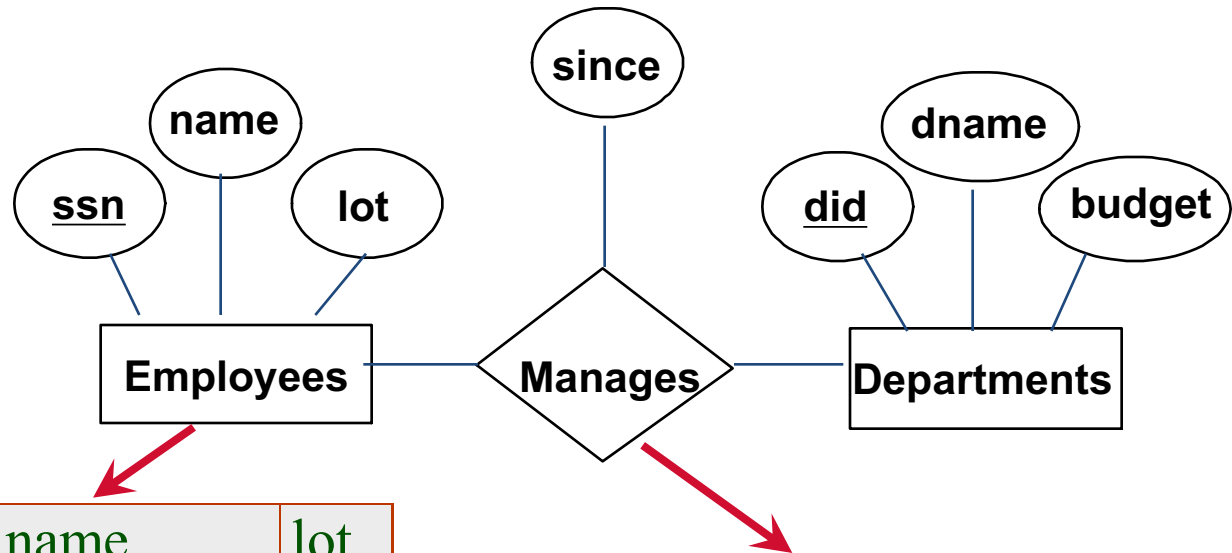


ssn	name	lot
123-22-3666	Attishoo	48
231-31-5368	Smiley	22
131-24-3650	Smethurst	35

```
CREATE TABLE Employees  
(ssn CHAR(11),  
name CHAR(20),  
lot INTEGER,  
PRIMARY KEY (ssn))
```

# Relationship Sets to Tables

Our favorite example:



<u>ssn</u>	name	lot
123-22-3666	Attishoo	48
231-31-5368	Smiley	22
131-24-3650	Smethurst	35

<u>ssn</u>	<u>did</u>	since
123-22-3666	51	1/1/91
123-22-3666	56	3/3/93
231-31-5368	51	2/2/92

# Relationship Sets to Tables

In translating a **many-to-many** relationship set to a relation, attributes of the relation must include:

Keys for each participating entity set (as foreign keys). This set of attributes forms a **superkey** for the relation.

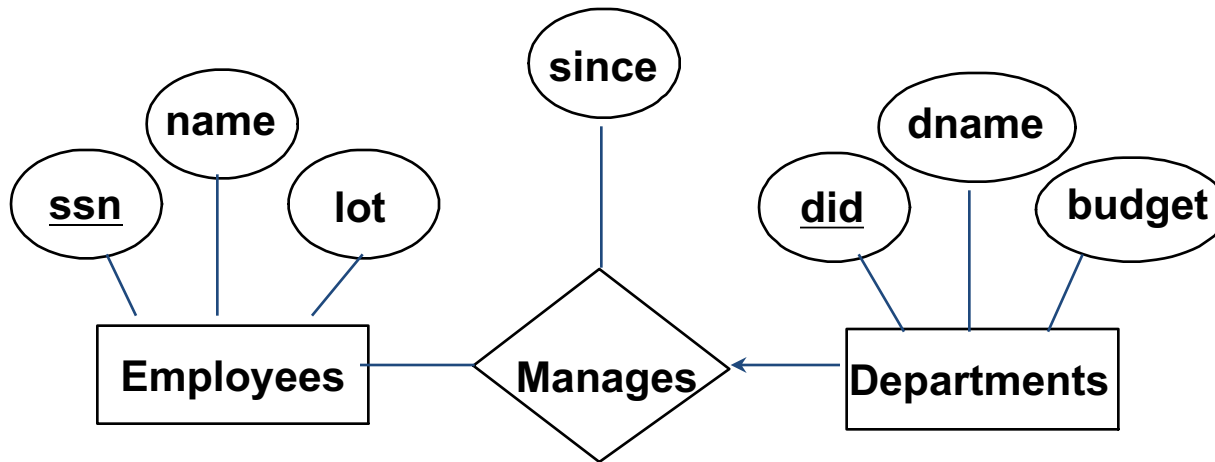
All descriptive attributes.

```
CREATE TABLE Manages(  
    ssn CHAR(11),  
    did INTEGER,  
    since DATE,  
    PRIMARY KEY (ssn, did),  
    FOREIGN KEY (ssn)  
        REFERENCES Employees,  
    FOREIGN KEY (did)  
        REFERENCES Departments)
```

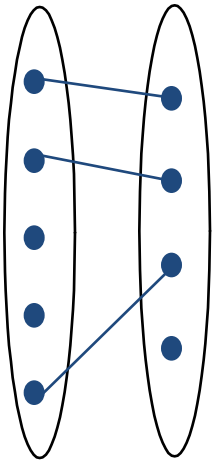
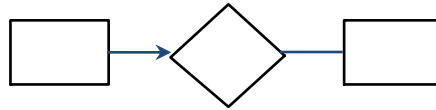
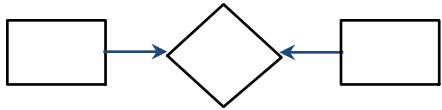
ssn	did	since
123-22-3666	51	1/1/91
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231-31-5368	51	2/2/92

# Review: Key Constraints in ER

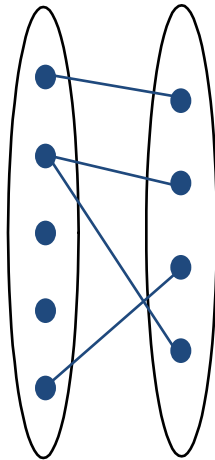
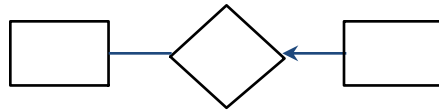
Each dept has at most one manager, according to the key constraint on Manages



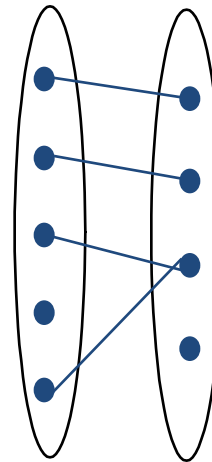
# Review: Key Constraints in ER



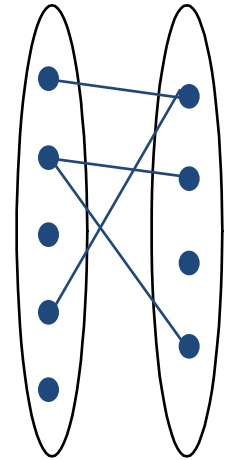
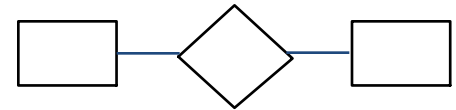
**1-to-1**



**1-to Many**

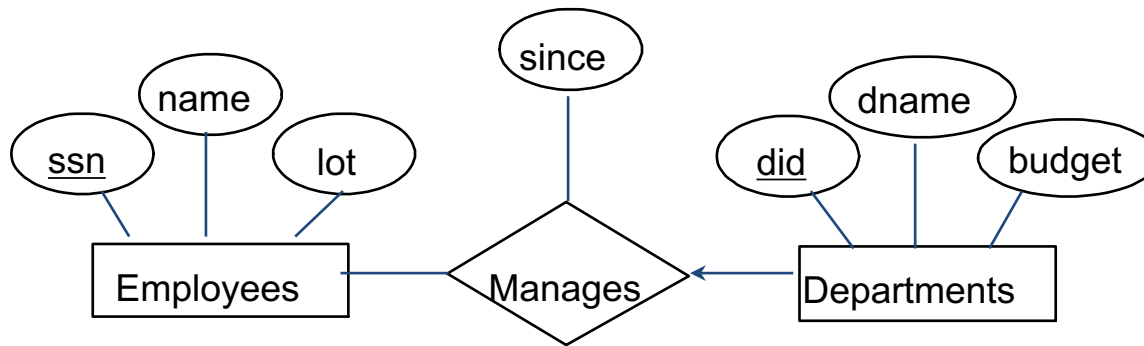


**Many-to-1**



**Many-to-Many**

# Translating ER with Key Constraints



since each department has a unique manager, we could instead combine Manages and Departments

```
CREATE TABLE Manages(  
  ssn CHAR(11),  
  did INTEGER,  
  since DATE,  
  PRIMARY KEY (did),  
  FOREIGN KEY (ssn) REFERENCES  
  Employees,  
  FOREIGN KEY (did) REFERENCES  
  Departments)
```

Vs.

```
CREATE TABLE Dept_Mgr(  
  did INTEGER,  
  dname CHAR(20),  
  budget REAL,  
  ssn CHAR(11),  
  since DATE,  
  PRIMARY KEY (did),  
  FOREIGN KEY (ssn)  
  REFERENCES Employees)
```

# What if the toy department has no manager (yet) ?



Can be NULL!



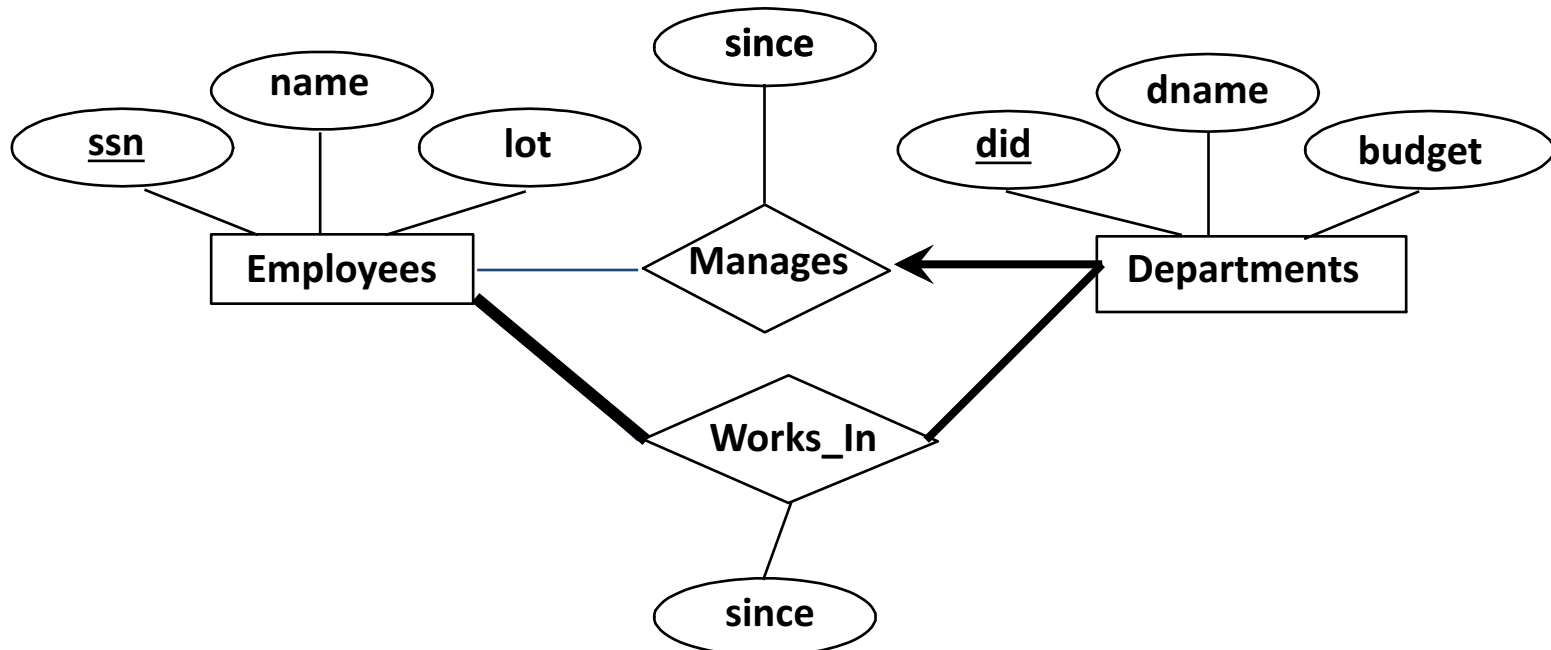
```
CREATE TABLE Dept_Mgr(  
  did INTEGER,  
  dname CHAR(20),  
  budget REAL,  
  ssn CHAR(11),  
  since DATE,  
  PRIMARY KEY (did),  
  FOREIGN KEY (ssn)  
  REFERENCES Employees)
```

# Review: Participation Constraints

does every employee work in a department?

If so, this is a participation constraint: the participation of Departments in Manages is said to be *total (vs. partial)*

Every *did* value in Departments table must appear in a row of the Manages table (with a non-null *ssn* value!)





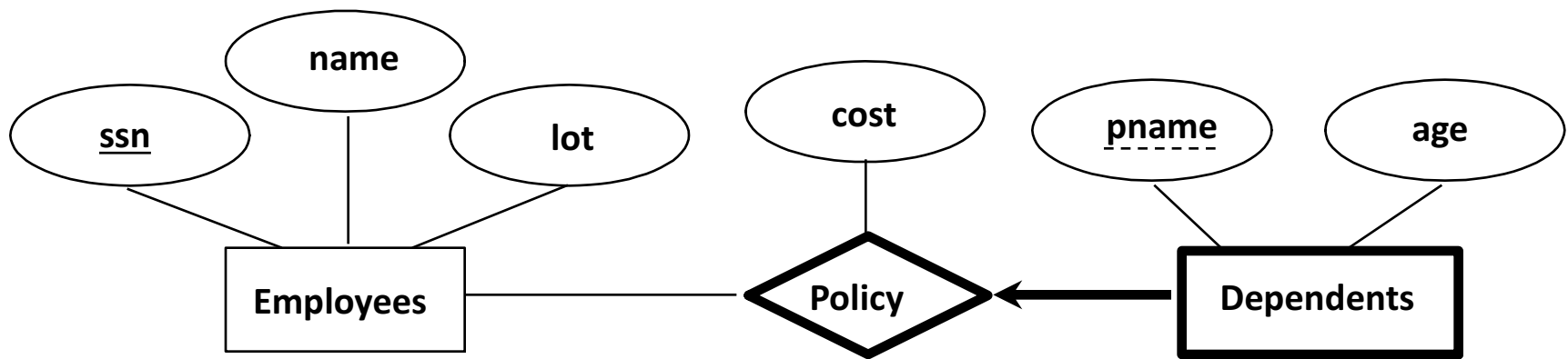
# Participation Constraints (PC) in SQL

PCs of one entity set in a binary relationship, yes!  
but little else (without resorting to CHECK constraints)

```
CREATE TABLE Dept_Mgr(  
  did INTEGER,  
  dname CHAR(20),  
  budget REAL,  
  ssn CHAR(11) NOT NULL,  
  since DATE,  
  PRIMARY KEY (did),  
  FOREIGN KEY (ssn) REFERENCES Employees,  
  ON DELETE NO ACTION)
```

# Review: Weak Entities

- A *weak entity* can be identified uniquely by the primary key of another (*owner*) entity (+ some of its attributes)
- Owner entity set and weak entity set must participate in a one-to-many relationship set (one owner, many weak entities)
  - Weak entity set must have total participation in this *identifying* relationship set



# Translating Weak Entity Sets

Weak entity set and identifying relationship set are translated into a single table.

When the owner entity is deleted, all owned weak entities must also be deleted.

```
CREATE TABLE Dep_Policy (  
  pname CHAR(20),  
  age INTEGER,  
  cost REAL,  
  ssn CHAR(11) NOT NULL,  
  PRIMARY KEY (pname, ssn),  
  FOREIGN KEY (ssn) REFERENCES Employees,  
  ON DELETE CASCADE)
```

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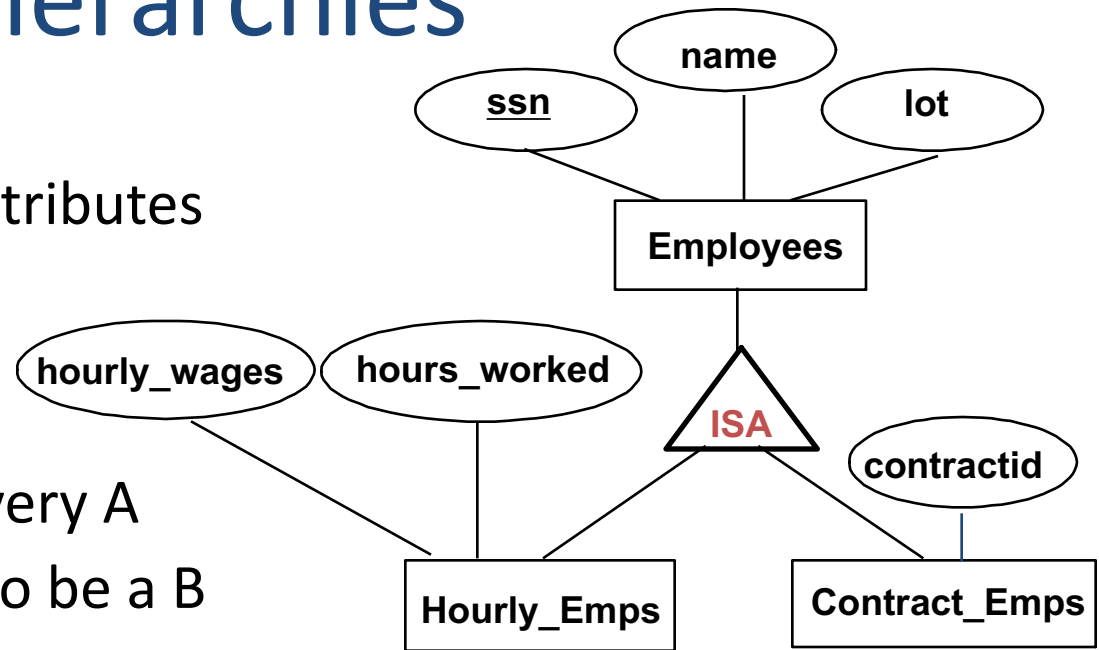
ER to Relational

**ISA to Relational**

# Review: ISA Hierarchies

As in C++, or other PLs, attributes are inherited.

If we declare A **ISA** B, every A entity is also considered to be a B entity.



*Overlap constraints:* Can Joe be an Hourly\_Emps as well as a Contract\_Emps entity? *(Allowed/disallowed)*

*Covering constraints:* Does every Employees entity also have to be an Hourly\_Emps or a Contract\_Emps entity? *(Yes/no)*

# Translating ISA Hierarchies to Relations

```
CREATE TABLE Employees (  
  ssn CHAR(11) NOT NULL,  
  name CHAR(20),  
  lot INTEGER,  
  PRIMARY KEY (ssn))
```

```
CREATE TABLE Contract_Emps  
(ssn CHAR(11) NOT NULL,  
  contractid INTEGER,  
  PRIMARY KEY (ssn),  
  FOREIGN KEY (ssn)  
    REFERENCES Employees))
```

```
CREATE TABLE Hourly_Emps (  
  ssn CHAR(11) NOT NULL,  
  hourly_wages REAL,  
  hours_worked REAL,  
  PRIMARY KEY (ssn),  
  FOREIGN KEY (ssn) REFERENCES Employees)
```

what should happen if I delete an entry from Employees?

can we use ON DELETE CASCADE?

how to access name and hours worked ?

Join!



# Alternative approach for ISA Hierarchies

```
CREATE TABLE Hourly_Emps (  
  ssn CHAR(11) NOT NULL,  
  name CHAR(20),  
  lot INTEGER,  
  hourly_wages REAL,  
  hours_worked REAL,  
  PRIMARY KEY (ssn))
```

```
CREATE TABLE Contract_Emps  
(ssn CHAR(11) NOT NULL,  
  name CHAR(20),  
  lot INTEGER,  
  contractid INTEGER,  
  PRIMARY KEY (ssn))
```

how to ensure that every employee is only in one of the two?



what about Employees that are neither?



what about querying for all employees?



Query 2 tables!

# Relational Model: Summary

tabular representation of data

simple & intuitive, currently the most widely used

*Integrity Constraints* can be specified based on app semantics & DBMS checks for violations

two important ICs: primary and foreign keys

in addition, we *always* have domain constraints

ER to Relational is (fairly) straightforward